

Probabilistic Graphical Models

Theory of Variational Inference: Inner and Outer Approximation



Junming Yin Lecture 15, March 4, 2013



Reading: W & J Book Chapters

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Roadmap



- Two families of approximate inference algorithms
 - Loopy belief propagation (sum-product)
 - Mean-field approximation
- Are there some connections of these two approaches?
- We will re-exam them from a unified point of view based on the variational principle:
 - Loop BP: outer approximation
 - Mean-field: inner approximation

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Variational Methods



- "Variational": fancy name for optimization-based formulations
 - i.e., represent the quantity of interest as the solution to an optimization problem
 - approximate the desired solution by relaxing/approximating the intractable optimization problem
- Examples:
 - Courant-Fischer for eigenvalues: $\lambda_{\max}(A) = \max_{\|x\|_2 = 1} x^T A x$
 - Linear system of equations: $Ax = b, A \succ 0, x^* = A^{-1}b$
 - variational formulation:

$$x^* = \arg\min_{x} \left\{ \frac{1}{2} x^T A x - b^T x \right\}$$

• for large system, apply conjugate gradient method

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Inference Problems in Graphical Models



• Undirected graphical model (MRF):

$$p(x) = \frac{1}{Z} \prod_{C \in \mathcal{C}} \psi_C(x_C)$$

- The quantities of interest:
 - ullet marginal distributions: $p(x_i) = \sum_{x_j, j
 eq i} p(x)$
 - ullet normalization constant (partition function): Z
- Question: how to represent these quantities in a variational form?
 - Use tools from (1) exponential families; (2) convex analysis

Exponential Families



· Canonical parameterization

$$p_{\theta}(x_1, \dots, x_m) = \exp \left\{ \theta^{\top} \phi(x) - A(\theta) \right\}$$

Canonical Parameters Sufficient Statistics Log partition Function

Log normalization constant:

$$A(\theta) = \log \int \exp\{\theta^T \phi(x)\} dx$$

- it is a convex function (Prop 3.1)
- Effective canonical parameters:

$$\Omega := \left\{ \theta \in \mathbb{R}^d | A(\theta) < +\infty \right\}$$

Graphical Models as Exponential Families



• Undirected graphical model (MRF):

$$p(\mathbf{x}; \theta) = \frac{1}{Z(\theta)} \prod_{C \in \mathcal{C}} \psi(\mathbf{x}_C; \theta_C)$$

MRF in an exponential form:

$$p(\mathbf{x}; \theta) = \exp \left\{ \sum_{C \in \mathcal{C}} \log \psi(\mathbf{x}_C; \theta_C) - \log Z(\theta) \right\}$$

• $\log \psi(\mathbf{x}_C; \theta_C)$ can be written in a *linear* form after some parameterization

Example: Gaussian MRF



- Consider a zero-mean multivariate Gaussian distribution that respects the Markov property of a graph
 - Hammersley-Clifford theorem states that the precision matrix $\Lambda=\Sigma^{-1}$ also respects the graph structure





• Gaussian MRF in the exponential form

$$p(\mathbf{x}) = \exp\left\{\frac{1}{2}\left\langle\Theta,\mathbf{x}\mathbf{x}^T\right\rangle - A(\Theta)\right\}, \text{where }\Theta = -\Lambda$$

 $\quad \text{Sufficient statistics are} \quad \{x_s^2, s \in V; x_s x_t, (s,t) \in E\}$

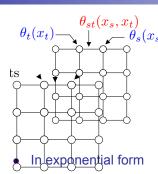
<u>Indicators:</u>

Parameters:

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Example: Discrete MRF





$$\mathbb{I}_{j}(x_{s}) = \begin{cases} 1 & \text{if } x_{s} = j \\ 0 & \text{otherwise} \end{cases}$$

$$\theta_s = \{\theta_{s;j}, j \in \mathcal{X}_s\}$$

$$\theta_{st} = \{\theta_{st;jk}, (j,k) \in \mathcal{X}_s \times \mathcal{X}_t\}$$

$$p(x;\theta) \propto \exp \left\{ \sum_{s \in V} \sum_{j} \theta_{s;j} \mathbb{I}_{j}(x_{s}) + \sum_{(s,t) \in E} \theta_{st;jk} \mathbb{I}_{j}(x_{s}) \mathbb{I}_{k}(x_{t}) \right\}$$

Why Exponential Families?



 Computing the expectation of sufficient statistics (mean parameters) given the canonical parameters yields the marginals

$$\mu_{s;j} = \mathbb{E}_p[\mathbb{I}_j(X_s)] = \mathbb{P}[X_s = j] \quad \forall j \in \mathcal{X}_s,$$

$$\mu_{st;jk} = \mathbb{E}_p[\mathbb{I}_{st;jk}(X_s, X_t)] = \mathbb{P}[X_s = j, X_t = k] \quad \forall (j,k) \in \mathcal{X}_s \in \mathcal{X}_t.$$

• Computing the normalizer yields the log partition function

$$\log Z(\theta) = A(\theta)$$

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Computing Mean Parameter: Bernoulli



• A single Bernoulli random variable

$$(X) \theta$$

$$p(x;\theta) = \exp\{\theta x - A(\theta)\}, x \in \{0,1\}, A(\theta) = \log(1 + e^{\theta})$$

• Inference = Computing the mean parameter

$$\mu(\theta) = \mathbb{E}_{\theta}[X] = 1 \cdot p(X = 1; \theta) + 0 \cdot p(X = 0; \theta) = \frac{e^{\theta}}{1 + e^{\theta}}$$

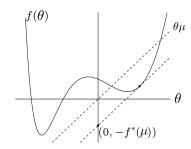
 Want to do it in a variational manner: cast the procedure of computing mean (summation) in an optimization-based formulation

Conjugate Dual Function



• Given any function $f(\theta)$, its conjugate dual function is:

$$f^*(\mu) = \sup_{\theta} \{ \langle \theta, \mu \rangle - f(\theta) \}$$



 Conjugate dual is always a convex function: point-wise supremum of a class of linear functions

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Dual of the Dual is the Original



• Under some technical condition on f (convex and lower semicontinuous), the dual of dual is itself:

$$f = (f^*)^*$$

$$f(\theta) = \sup_{\mu} \left\{ \langle \theta, \mu \rangle - f^*(\mu) \right\}$$

• For log partition function

$$A(\theta) = \sup_{\mu} \{ \langle \theta, \mu \rangle - A^*(\mu) \}, \quad \theta \in \Omega$$

ullet The dual variable μ has a natural interpretation as the mean parameters

Computing Mean Parameter: Bernoulli



- The conjugate $A^*(\mu) := \sup_{\theta \in \mathbb{R}} \{\mu\theta \log[1 + \exp(\theta)]\}$ Stationary condition $\mu = \frac{e^{\theta}}{1 + e^{\theta}}$ $(\mu = \nabla A(\theta))$
- If $\mu \in (0,1)$, $\theta(\mu) = \log\left(\frac{\mu}{1-\mu}\right)$, $A^*(\mu) = \mu \log(\mu) + (1-\mu) \log(1-\mu)$
- $\bullet \quad \text{We have} \quad A^*(\mu) = \begin{cases} \mu \log \mu + (1-\mu) \log (1-\mu) & \text{if } \mu \in [0,1] \\ +\infty & \text{therwise.} \end{cases}$ $\bullet \quad \text{The variational form:} \quad A(\theta) = \max_{\mu \in [0,1]} \left\{ \mu \cdot \theta + A^*(\mu) \right\}.$
- \bullet The optimum is achieved at $\,\mu(\theta)=\frac{e^{\theta}}{1+e^{\theta}}\,$. This is the mean!

Remark



- The last few identities are not coincidental but rely on a deep theory in general exponential family.
 - The dual function is the negative entropy function
 - The mean parameter is restricted
 - Solving the optimization returns the mean parameter and log partition function
- Next step: develop this framework for general exponential families/graphical models.
- However,
 - Computing the conjugate dual (entropy) is in general intractable
 - The constrain set of mean parameter is hard to characterize
 - Hence we need approximation

Computation of Conjugate Dual



• Given an exponential family

$$p(x_1, \dots, x_m; \theta) = \exp \left\{ \sum_{i=1}^d \theta_i \phi_i(x) - A(\theta) \right\}$$

The dual function

$$A^*(\mu) := \sup_{\theta \in \Omega} \{ \langle \mu, \theta \rangle - A(\theta) \}$$

- The stationary condition: $\mu \nabla A(\theta) = 0$
- Derivatives of A yields mean parameters

$$\frac{\partial A}{\partial \theta_i}(\theta) = \mathbb{E}_{\theta}[\phi_i(X)] = \int \phi_i(x)p(x;\theta) dx$$

- The stationary condition becomes $\mu = \mathbb{E}_{\theta}[\phi(X)]$
- Question: for which $\mu \in \mathbb{R}^d$ does it have a solution $\theta(\mu)$?

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Computation of Conjugate Dual



- Let's assume there is a solution $\theta(\mu)$ such that $\mu = \mathbb{E}_{\theta(u)}[\phi(X)]$
- The dual has the form

$$A^{*}(\mu) = \langle \theta(\mu), \mu \rangle - A(\theta(\mu))$$

$$= \mathbb{E}_{\theta(\mu)} \left[\langle \theta(\mu), \frac{\phi(X)}{\rho(X)} \rangle - A(\theta(\mu)) \right]$$

$$= \mathbb{E}_{\theta(\mu)} \left[\log p(X; \theta(\mu)) \right]$$

• The entropy is defined as

$$H(p(x)) = -\int p(x) \log p(x) dx$$

• So the dual is $A^*(\mu) = -H(p(x; \theta(\mu)))$ when there is a solution $\theta(\mu)$

Complexity of Computing Conjugate Dual



• The dual function is implicitly defined:

$$\mu \qquad - \blacktriangleright \boxed{(\nabla A)^{-1} \qquad \stackrel{\theta(\mu)}{\blacktriangleright} \qquad -H(p_{\theta(\mu)}) \qquad \stackrel{}{\blacktriangleright} \quad A^*(\mu)$$

- Solving the inverse mapping $\mu = \mathbb{E}_{\theta}[\phi(X)]$ for canonical parameters $\theta(\mu)$ is nontrivial
- Evaluating the negative entropy requires high-dimensional integration (summation)
- Question: for which $\mu \in \mathbb{R}^d$ does it have a solution $\theta(\mu)$? i.e., the domain of $A^*(\mu)$.
 - the ones in marginal polytope!

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Marginal Polytope



• For any distribution p(x) and a set of sufficient statistics $\phi'(x)$, define a vector of mean parameters

$$\mu_i = \mathbb{E}_p[\phi_i(X)] = \int \phi_i(x)p(x) dx$$

- p(x) is not necessarily an exponential family
- The set of all realizable mean parameters

$$\mathcal{M} := \{ \mu \in \mathbb{R}^d \mid \exists p \text{ s.t. } \mathbb{E}_p[\phi(X)] = \mu \}.$$

- It is a convex set
- For discrete exponential families, this is called marginal polytope

Convex Polytope



Convex hull representation

$$\begin{split} \mathcal{M} &= \Big\{ \mu \in \mathbb{R}^d | \sum_{x \in \mathcal{X}^m} \phi(x) p(x) = \mu, \text{ for some } p(x) \geq 0, \sum_{x \in \mathcal{X}^m} p(x) = 1 \Big\} \\ &\triangleq \text{conv} \Big\{ \phi(x), x \in \mathcal{X}^m \Big\} \end{split}$$

- Half-plane representation
 - Minkowski-Weyl Theorem: Payameterization and Juferencyt Problems be 55 characterized by a finite collection of linear inequality constraints

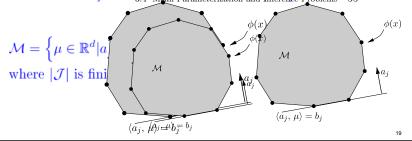


Fig. 3.5 Generic illustration of \mathcal{M} for a discrete random variable with $|\mathcal{X}^m|$ finite. In this Fig. 3.5 Generic illustration one \mathcal{M} for a discrete random variable with $|\mathcal{X}^m|$ finite \mathcal{M}^m , this this can be sufficiently finite \mathcal{M}^m . case, the spy M in a convex vpolytope rectangues porting tenths convex thull as M and MBy the Minkowski-More threshes partition of a finite unitable of half-spaces, each of the form $\{\mu \in \mathbb{R}^d \mid \langle a_j, \mu \rangle \geq b_j \}$ for some pair $(a_j,b_j) \in \mathbb{R}^d \times \mathbb{R}.$ To make these ideas in the scherece; consider the schriftest on ideas in the simplest nontrivial

case: namely, a pair of variablely (XPax, of variables (XL), Xonsisting the graph consisting of Example 3.8 (Ising Wienin Parameters): which is supply to the set M is a polytope in ple 3.1, the supply supply the set M is a polytope in ple 3.1, the supply supply

Example of the substitution of the substituti

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or the $cut\ polytope_{\mathbb{R}}$ 69ull 87pm of Eight Hebelipuus A für ihm smoini sen awan krimelenodel with two variables $(X_1,X_2)\in\{0,1\}^n$ The fold with final failure and the constraints satisfy the constraints supply the supply of the constraints satisfy the constraints of the

These four constraints provide an alternative characterization of the 3D polytope illustrated in Figure 3.6.

Marginal Polytope for General Graphs

- Still doable for connected binary graphs with 3 nodes: 16 constraints
- For tree graphical models, the number of half-planes (facet complexity) grows only *linearly* in the graph size
- General graphs?
 - extremely hard to characterize the marginal polytope



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Variational Principle (Theorem 3.4)



• The dual function takes the form

$$A^*(\mu) = \begin{cases} -H(p_{\theta(\mu)}) & \text{if } \mu \in \mathcal{M}^{\circ} \\ +\infty & \text{if } \mu \notin \overline{\mathcal{M}}. \end{cases}$$

- $\bullet \quad \theta(\mu) \ \ \text{satisfies} \ \ \mu = \mathbb{E}_{\theta(u)}[\phi(X)]$
- The log partition function has the variational form

$$A(\theta) = \sup_{\mu \in \mathcal{M}} \{ \theta^T \mu - A^*(\mu) \}$$

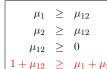
• For all $\theta \in \Omega$, the above optimization problem is attained uniquely at $\mu(\theta) \in \mathcal{M}^o$ that satisfies

$$\mu(\theta) = \mathbb{E}_{\theta}[\phi(X)]$$

Example: Two-node Ising Model



- The distribution $p(x;\theta) \propto \exp\{\theta_1 x_1 + \theta_2 x_2 + \theta_{12} x_{12}\}$
 - Sufficient statistics $\phi(x) = \{x_1, x_2, x_1x_2\}$



- The marginal polytope is characterized by
- The dual has an explicit form

$$A^*(\mu) = \mu_{12} \log \mu_{12} + (\mu_1 - \mu_{12}) \log(\mu_1 - \mu_{12}) + (\mu_2 - \mu_{12}) \log(\mu_2 - \mu_{12}) + (1 + \mu_{12} - \mu_1 - \mu_2) \log(1 + \mu_{12} - \mu_1 - \mu_2)$$

- The variational problem $A(\theta) = \max_{\{\mu_1, \mu_2, \mu_{12}\} \in \mathcal{M}} \{\theta_1 \mu_1 + \theta_2 \mu_2 + \theta_{12} \mu_{12} A^*(\mu)\}$
- The optimum is attained at

$$\mu_1(\theta) = \frac{\exp\{\theta_1\} + \exp\{\theta_1 + \theta_2 + \theta_{12}\}}{1 + \exp\{\theta_1\} + \exp\{\theta_2\} + \exp\{\theta_1 + \theta_2 + \theta_{12}\}}$$

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Variational Principle



• Exact variational formulation

$$A(\theta) = \sup_{\mu \in \mathcal{M}} \{ \theta^T \mu - A^*(\mu) \}$$

- \mathcal{M} : the marginal polytope, difficult to characterize
- A^* : the negative entropy function, no explicit form
- Mean field method: non-convex inner bound and exact form of entropy
- Bethe approximation and loopy belief propagation: polyhedral outer bound and non-convex Bethe approximation



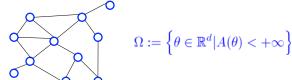
Mean Field Approximation

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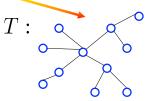
Tractable Subgraphs



- Definition: A subgraph F of the graph G is *tractable* if it is feasible to perform exact inference
- Example:







$$\Omega(F_0) := \{\theta \in \Omega | \theta_{(s,t)} = 0, \forall (s,t) \in E\} \quad \Omega(T) := \{\theta \in \Omega | \theta_{(s,t)} = 0 \ \forall (s,t) \notin E(T)\}$$

Mean Field Methods



ullet For an exponential family with sufficient statistics ϕ defined on graph G, the set of realizable mean parameter set

$$\mathcal{M}(G;\phi) := \{ \mu \in \mathbb{R}^d \mid \exists p \text{ s.t. } \mathbb{E}_p[\phi(X)] = \mu \}$$

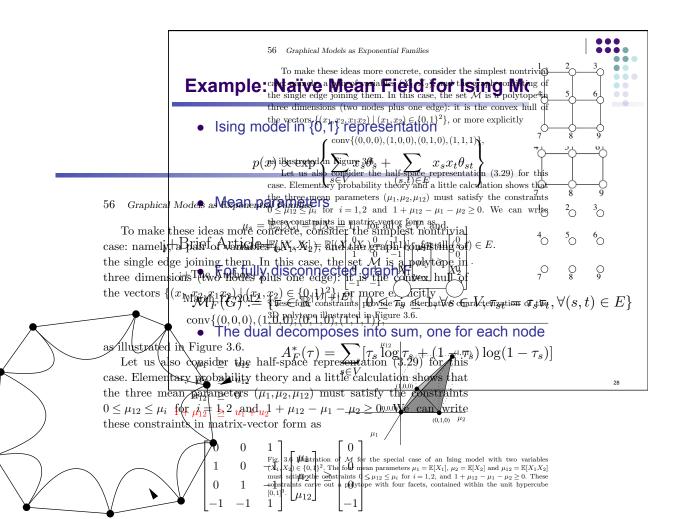
 For a given tractable subgraph F, a subset of mean parameters of interest

$$\mathcal{M}(F;\phi) := \{ \tau \in \mathbb{R}^d \mid \tau = \mathbb{E}_{\theta}[\phi(X)] \text{ for some } \theta \in \Omega(F) \}$$

- Inner approximation $\mathcal{M}(F;\phi)^o \subseteq \mathcal{M}(G;\phi)^o$
- Mean field solves the relaxed problem

$$\max_{\tau \in \mathcal{M}_F(G)} \{ \langle \tau, \theta \rangle - A_F^*(\tau) \}$$

• $A_F^* = A^* \big|_{\mathcal{M}_F(G)}$ is the exact dual function restricted to $\mathcal{M}_F(G)$



Example: Naïve Mean Field for Ising Model



Mean field problem

$$A(\theta) \ge \max_{(\tau_1, \dots, \tau_m) \in [0, 1]^m} \left\{ \sum_{s \in V} \theta_s \tau_s + \sum_{(s, t) \in E} \theta_{st} \tau_s \tau_t - A_F^*(\tau) \right\}$$

- The same objective function as in free energy based approach
- The naïve mean field update equations

$$\tau_s \leftarrow \sigma \left(\theta_s + \sum_{t \in N(s)} \theta_s \tau_t\right)$$

Also yields lower bound on log partition function

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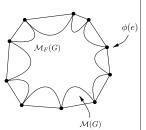
Geometry of Mean Field



- Mean field optimization is always non-convex for any exponential family in which the state space \mathcal{X}^m is finite
- Recall the marginal polytope is a convex hull

$$\mathcal{M}(G) = \operatorname{conv}\{\phi(e); e \in \mathcal{X}^m\}$$

- $\mathcal{M}_F(G)$ contains all the extreme points
 - If it is a strict subset, then it must be non-convex



• Example: two-node Ising model

$$\mathcal{M}_F(G) = \{0 \le \tau_1 \le 1, 0 \le \tau_2 \le 1, \tau_{12} = \tau_1 \tau_2\}$$

• It has a parabolic cross section along $\, au_1 = au_2\,$, hence non-convex



Bethe Approximation and Sum-Product

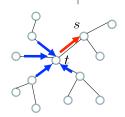
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Sum-Product/Belief Propagation Algorithm



• Message passing rule:

$$M_{ts}(x_s) \leftarrow \kappa \sum_{x_t'} \left\{ \psi_{st}(x_s, x_t') \psi_t(x_t') \prod_{u \in N(t)/s} M_{ut}(x_t') \right\}$$



• Marginals:

$$\mu_s(x_s) = \kappa \, \psi_s(x_s) \prod_{t \in N(s)} M_{ts}^*(x_s)$$

- Exact for trees, but approximate for loopy graphs (so called loopy belief propagation)
- Question:
 - How is the algorithm on trees related to variational principle?
 - What is the algorithm doing for graphs with cycles?

Tree Graphical Models



- Discrete variables $X_s \in \{0, 1, \dots, m_s 1\}$ on a tree T = (V, E)
- Exponential representation of distribution:

$$p(\mathbf{x};\theta) \propto \exp\big\{\sum_{s \in V} \theta_s(x_s) + \sum_{(s,t) \in E} \theta_{st}(x_s,x_t)\big\}$$
 where $\theta_s(x_s) := \sum_{j \in \mathcal{X}_s} \theta_{s;j} \mathbb{I}_j(x_s)$ (and similarly for $\theta_{st}(x_s,x_t)$)

Mean partiers are marginal probabilities:

$$\mu_{s;j} = \mathbb{E}_{p}[\mathbb{I}_{j}(X_{s})] = \mathbb{P}[X_{s} = j] \quad \forall j \in \mathcal{X}_{s}, \quad \mu_{s}(x_{s}) = \sum_{j \in \mathcal{X}_{s}} \mu_{s;j} \mathbb{I}_{j}(x_{s}) = \mathbb{P}(X_{s} = x_{s})$$

$$\mu_{st;jk} = \mathbb{E}_{p}[\mathbb{I}_{st;jk}(X_{s}, X_{t})] = \mathbb{P}[X_{s} = j, X_{t} = k] \quad \forall (j,k) \in \mathcal{X}_{s} \in \mathcal{X}_{t}.$$

$$\mu_{st;jk}\mathbb{I}_{jk}(x_{s}, x_{t}) = \mathbb{P}(X_{s} = x_{s}, X_{t} = x_{t})$$

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Marginal Polytope for Trees



• Recall marginal polytope for general graphs

$$\mathcal{M}(G) = \{ \mu \in \mathbb{R}^d \mid \exists p \text{ with marginals } \mu_{s;j}, \mu_{st;jk} \}$$

• By junction tree theorem (see Prop. 2.1 & Prop. 4.1)

$$\mathcal{M}(T) = \left\{ \mu \ge 0 \mid \sum_{x_s} \mu_s(x_s) = 1, \sum_{x_t} \mu_{st}(x_s, x_t) = \mu_s(x_s) \right\}$$

• In particular, if $\mu \in \mathcal{M}(T)$, then

$$p_{\mu}(x) := \prod_{s \in V} \mu_s(x_s) \prod_{(s,t) \in E} \frac{\mu_{st}(x_s, x_t)}{\mu_s(x_s)\mu_t(x_t)}.$$

has the corresponding marginals





• For trees, the entropy decomposes as

$$\begin{split} H(p(x;\mu)) &= -\sum_{x} p(x;\mu) \log p(x;\mu) \\ &= \sum_{s \in V} \left(-\sum_{\underline{x_s}} \mu_s(x_s) \log \mu_s(x_s) \right) - \\ &- \sum_{(s,t) \in E} \left(\sum_{\underline{x_s},x_t} \mu_{st}(x_s,x_t) \log \frac{\mu_{st}(x_s,x_t)}{\mu_s(x_s)\mu_t(x_t)} \right) \\ &= \sum_{s \in V} H_s(\mu_s) - \sum_{(s,t) \in E} I_{st}(\mu_{st}) \end{split}$$

• The dual function has an explicit form $A^*(\mu) = -H(p(x;\mu))$

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Exact Variational Principle for Trees



Variational formulation

$$A(\theta) = \max_{\mu \in \mathcal{M}(T)} \left\{ \langle \theta, \mu \rangle + \sum_{s \in V} H_s(\mu_s) - \sum_{(s,t) \in E} I_{st}(\mu_{st}) \right\}$$

- Assign Lagrange multiplier λ_{ss} for the normalization constraint $C_{ss}(\mu) := 1 \sum_{x_s} \mu_s(x_s) = 0$; and $\lambda_{ts}(x_s)$ for each marginalization constraint $C_{ts}(x_s;\mu) := \mu_s(x_s) \sum_{x_t} \mu_{st}(x_s,x_t) = 0$
- The Lagrangian has the form

$$\mathcal{L}(\mu, \lambda) = \langle \theta, \mu \rangle + \sum_{s \in V} H_s(\mu_s) - \sum_{(s,t) \in E} I_{st}(\mu_{st}) + \sum_{s \in V} \lambda_{ss} C_{ss}(\mu)$$
$$+ \sum_{(s,t) \in E} \left[\sum_{x_t} \lambda_{st}(x_t) C_{st}(x_t) + \sum_{x_s} \lambda_{ts}(x_s) C_{ts}(x_s) \right]$$

Lagrangian Derivation



• Taking the derivatives of the Lagrangian w.r.t. μ_s and μ_{st}

$$\frac{\partial \mathcal{L}}{\partial \mu_s(x_s)} = \theta_s(x_s) - \log \mu_s(x_s) + \sum_{t \in \mathcal{N}(s)} \lambda_{ts}(x_s) + C$$

$$\frac{\partial \mathcal{L}}{\partial \mu_{st}(x_s, x_t)} = \theta_{st}(x_s, x_t) - \log \frac{\mu_{st}(x_s, x_t)}{\mu_s(x_s)\mu_t(x_t)} - \lambda_{ts}(x_s) - \lambda_{st}(x_t) + C'$$

· Setting them to zeros yields

$$\mu_s(x_s) \propto \exp\{\theta_s(x_s)\} \prod_{t \in \mathcal{N}(s)} \underbrace{\exp\{\lambda_{ts}(x_s)\}}_{M_{ts}(x_s)}$$

$$\mu_s(x_s, x_t) \propto \exp\{\theta_s(x_s) + \theta_t(x_t) + \theta_{st}(x_s, x_t)\} \times \prod_{u \in \mathcal{N}(s) \setminus t} \exp\{\lambda_{us}(x_s)\} \prod_{v \in \mathcal{N}(t) \setminus s} \exp\{\lambda_{vt}(x_t)\}$$

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Lagrangian Derivation (continued)



Adjusting the Lagrange multipliers or messages to enforce

$$C_{ts}(x_s; \mu) := \mu_s(x_s) - \sum_{x_t} \mu_{st}(x_s, x_t) = 0$$

yields

$$M_{ts}(x_s) \leftarrow \sum_{x_t} \exp \left\{ \theta_t(x_t) + \theta_{st}(x_s, x_t) \right\} \prod_{u \in \mathcal{N}(t) \setminus s} M_{ut}(x_t)$$

 Conclusion: the message passing updates are a Lagrange method to solve the stationary condition of the variational formulation

BP on Arbitrary Graphs



Two main difficulties of the variational formulation

$$A(\theta) = \sup_{\mu \in \mathcal{M}} \{ \theta^T \mu - A^*(\mu) \}$$

 The marginal polytope M is hard to characterize, so let's use the treebased outer bound

$$\mathbb{L}(G) = \left\{ \tau \ge 0 \mid \sum_{x_s} \tau_s(x_s) = 1, \sum_{x_t} \tau_{st}(x_s, x_t) = \tau_s(x_s) \right\}$$

These locally consistent vectors τ are called pseudo-marginals.

• Exact entropy $-A^*(\mu)$ lacks explicit form, so let's approximate it by the exact expression for trees

$$-A^*(\tau) \approx H_{\text{Bethe}}(\tau) := \sum_{s \in V} H_s(\tau_s) - \sum_{(s,t) \in E} I_{st}(\tau_{st}).$$

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Bethe Variational Problem (BVP)



 Combining these two ingredient leads to the Bethe variational problem (BVP):

$$\max_{\tau \in \mathbb{L}(G)} \left\{ \langle \theta, \tau \rangle + \sum_{s \in V} H_s(\tau_s) - \sum_{(s,t) \in E} I_{st}(\tau_{st}) \right\}.$$

- A simple structured problem (differentiable & constraint set is a simple convex polytope)
- Loopy BP can be derived as am iterative method for solving a Lagrangian formulation of the BVP (Theorem 4.2); similar proof as for tree graphs

Geometry of BP

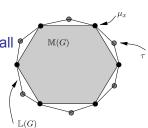


 $\begin{bmatrix} 0.4 & 0.1 \\ 0.1 & 0.4 \end{bmatrix}$

- Consider the following assignment of pseudo-marginals
 - Can easily verify $\tau \in \mathbb{L}(G)$
 - However, $\tau \notin \mathcal{M}(G)$ (need a bit more work)



- For any graph, $\mathcal{M}(G) \subseteq \mathbb{L}(G)$
- Equality holds if and only if the graph is a tree



- Question: does solution to the BVP ever fall into the gap?
 - Yes, for any element of outer bound L(G), it is possible to construct a distribution with it as a BP fixed point (Wainwright et. al. 2003)

Inexactness of Bethe Entropy Approximation



Consider a fully connected graph with

$$\mu_s(x_s) = \begin{bmatrix} 0.5 & 0.5 \end{bmatrix} \text{ for } s = 1, 2, 3, 4$$

$$\mu_{st}(x_s, x_t) = \begin{bmatrix} 0.5 & 0 \\ 0 & 0.5 \end{bmatrix} \quad \forall (s, t) \in E.$$



- It is globally valid: $\tau \in \mathcal{M}(G)$; realized by the distribution that places mass 1/2 on each of configuration (0,0,0,0) and (1,1,1,1)
- $H_{\text{Bethe}}(\mu) = 4\log 2 6\log 2 = -2\log 2 < 0$,
- $-A^*(\mu) = \log 2 > 0$.

Remark



- This connection provides a principled basis for applying the sum-product algorithm for loopy graphs
- However,
 - Although there is always a fixed point of loopy BP, there is no guarantees on the convergence of the algorithm on loopy graphs
 - The Bethe variational problem is usually non-convex. Therefore, there
 are no guarantees on the global optimum
 - ullet Generally, no guarantees that $A_{
 m Bethe}(heta)$ is a lower bound of A(heta)
- Nevertheless,
 - The connection and understanding suggest a number of avenues for improving upon the ordinary sum-product algorithm, via progressively better approximations to the entropy function and outer bounds on the marginal polytope (Kikuchi clustering)

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Summary



- Variational methods in general turn inference into an optimization problem via exponential families and convex duality
- The exact variational principle is intractable to solve; there are two distinct components for approximations:
 - Either inner or outer bound to the marginal polytope
 - Various approximation to the entropy function
- Mean field: non-convex inner bound and exact form of entropy
- BP: polyhedral outer bound and non-convex Bethe approximation
- <u>Kikuchi and variants</u>: tighter polyhedral outer bounds and better entropy approximations (Yedidia et. al. 2002)